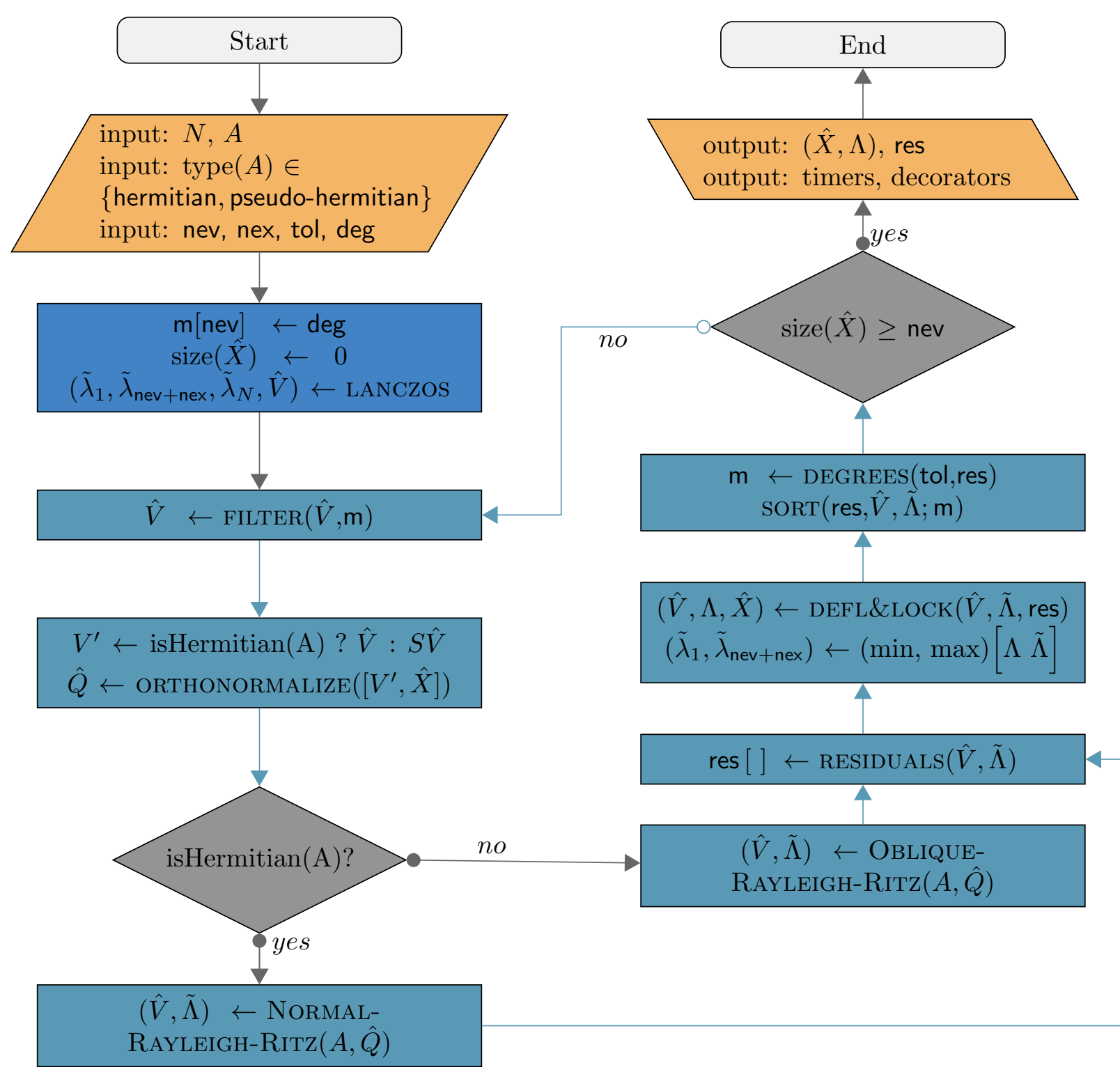


Chebyshev Accelerated Subspace Eigensolvers for Pseudo-hermitian Hamiltonians

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ChASE Algorithm



GitHub: <https://github.com/ChASE-library/ChASE>
License: Open source – BSD 3-Clause
Latest release: v1.7.0-rc2 – January, 2026
Docs: chase-library.github.io/ChASE/quick-start.html
Easy-to-integrate: ready-to-use C++ to Fortran interface
accessible via ELSI (Electronic Structure Infrastructure)

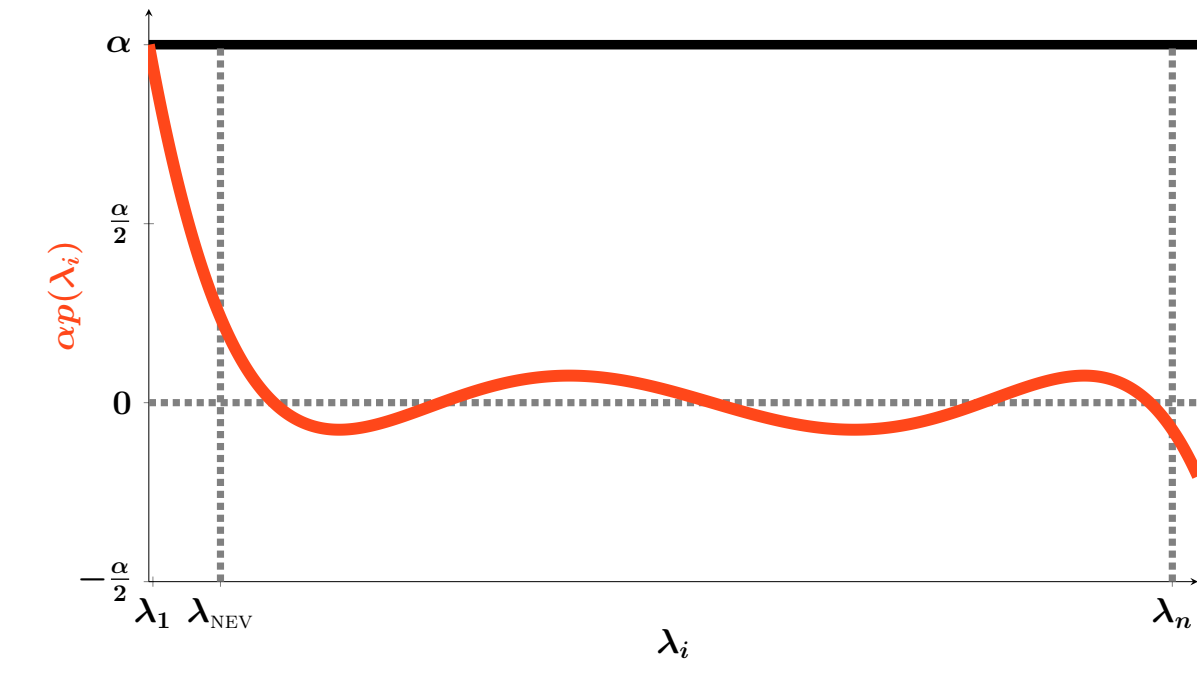


Figure: Illustration of the filter on a uniformly distributed random vector

Problem Definition

The Hamiltonian H derived by the BSE and by the DFT is of the form

$$H := \begin{bmatrix} A & B \\ -B & -A \end{bmatrix} \quad \text{with} \quad A = A^* \quad \text{and} \quad B = B^T.$$

The matrix H is said “pseudo-hermitian” as it satisfies

$$SH = H^*S \quad \text{with} \quad S := \begin{bmatrix} I & 0 \\ 0 & -I \end{bmatrix}.$$

This poster focuses on the most common case in which the eigenproblem

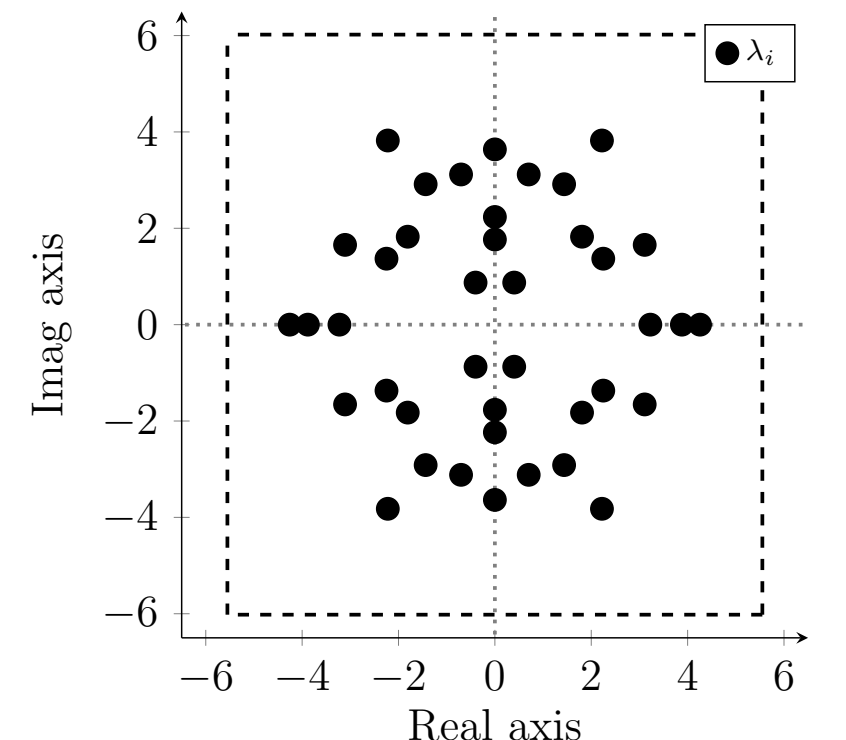
$$Hv = \lambda v \Leftrightarrow SHv = \lambda Sv \quad \text{with} \quad v \in \mathbb{C}^n \quad \text{and} \quad \lambda \in \mathbb{R}$$

is “definite”, meaning that SH is hermitian positive definite (HPD), i.e.,

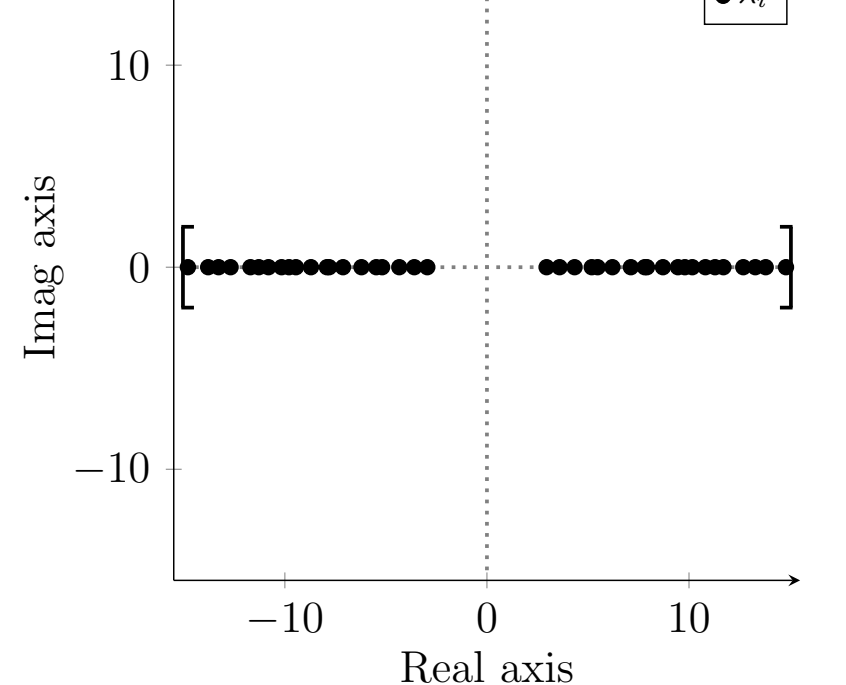
$$SH = \begin{bmatrix} A & B \\ B & -A \end{bmatrix} \succ 0.$$

Summary of the key properties :

- The spectrum is real, i.e., $\lambda(H) \in \mathbb{R}$.
- Eigenvalues appear in positive-negative pairs, such that $Hv = \lambda v \Leftrightarrow HKv = -\lambda Kv$ with $K := \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix}$.
- The left and right eigenvectors are related such that $U = SV$.
- The set of right eigenvectors is not orthonormal, but satisfies $V^*SV = D$ with D a non-singular real diagonal matrix.



(a) $\lambda_i(H) - SH$ is indefinite



(b) $\lambda_i(H) - SH$ is HPD

Major Algorithmic Innovations

Rayleigh-Ritz for pseudo-hermitian Hamiltonian

Orthonormal vs. Oblique Rayleigh-Ritz

- Hermitian case :
- Orthonormal Projection : $\tilde{y}^* (A\tilde{x} - \tilde{\omega}\tilde{x}) = 0, \quad \forall \tilde{y} \in \text{range}(Q).$
 - Quadratic Convergence : $|\omega - \tilde{\omega}| \leq |\sigma_v|^2 \|A - \omega I\|_2,$
 - Single basis : $\lim_{Q \rightarrow V} Q^*AQ = \Omega.$
- Non-hermitian case :
- Oblique Projection : $\tilde{u}^* (H\tilde{v} - \tilde{\lambda}\tilde{v}) = 0, \quad \text{with} \quad \tilde{v} \in \text{range}(Q) \quad \text{and} \quad \forall \tilde{u} \in \text{range}(Q_L).$
 - Non-immediate Quadratic Convergence : $|\lambda - \tilde{\lambda}| \leq |\tilde{\delta}|^{-1} \cdot \|\tilde{\sigma}_u\sigma_v\| \cdot \|H - \lambda I\|_2, \quad \tilde{\delta} := \tilde{u}^*\tilde{v}$
 - Dual basis, such that $Q_L^*Q = I:$ $\lim_{Q \rightarrow V, Q_L \rightarrow U} Q_L^*HQ = \Lambda.$

General Form of Dual Basis

Let M be a non-singular matrix. A dual basis Q_L is defined as

$$Q_L := [SQ - Q(Q^*SQ - M)]M^{-1}.$$

Since Q is orthonormal ($Q^*Q = I$), we have

$$\begin{aligned} Q_L^*Q &= ([SQ - Q(Q^*SQ - M)]M^{-1})^*Q \\ &= (M^{-1})^*(Q^*SQ - Q^*S^*Q + M^*) \\ &= (M^{-1})^*M^* = I. \end{aligned}$$

The Rayleigh-Quotient $G := Q_L^*HQ$ is not hermitian.

Flip-sign functions that emulate the multiplication by S

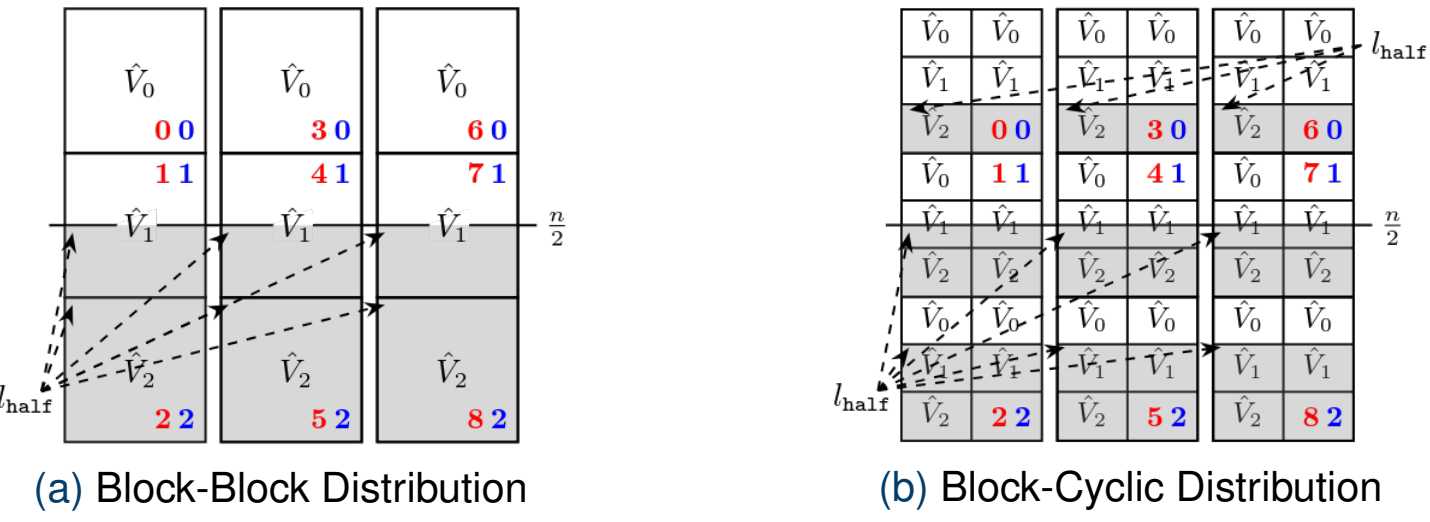
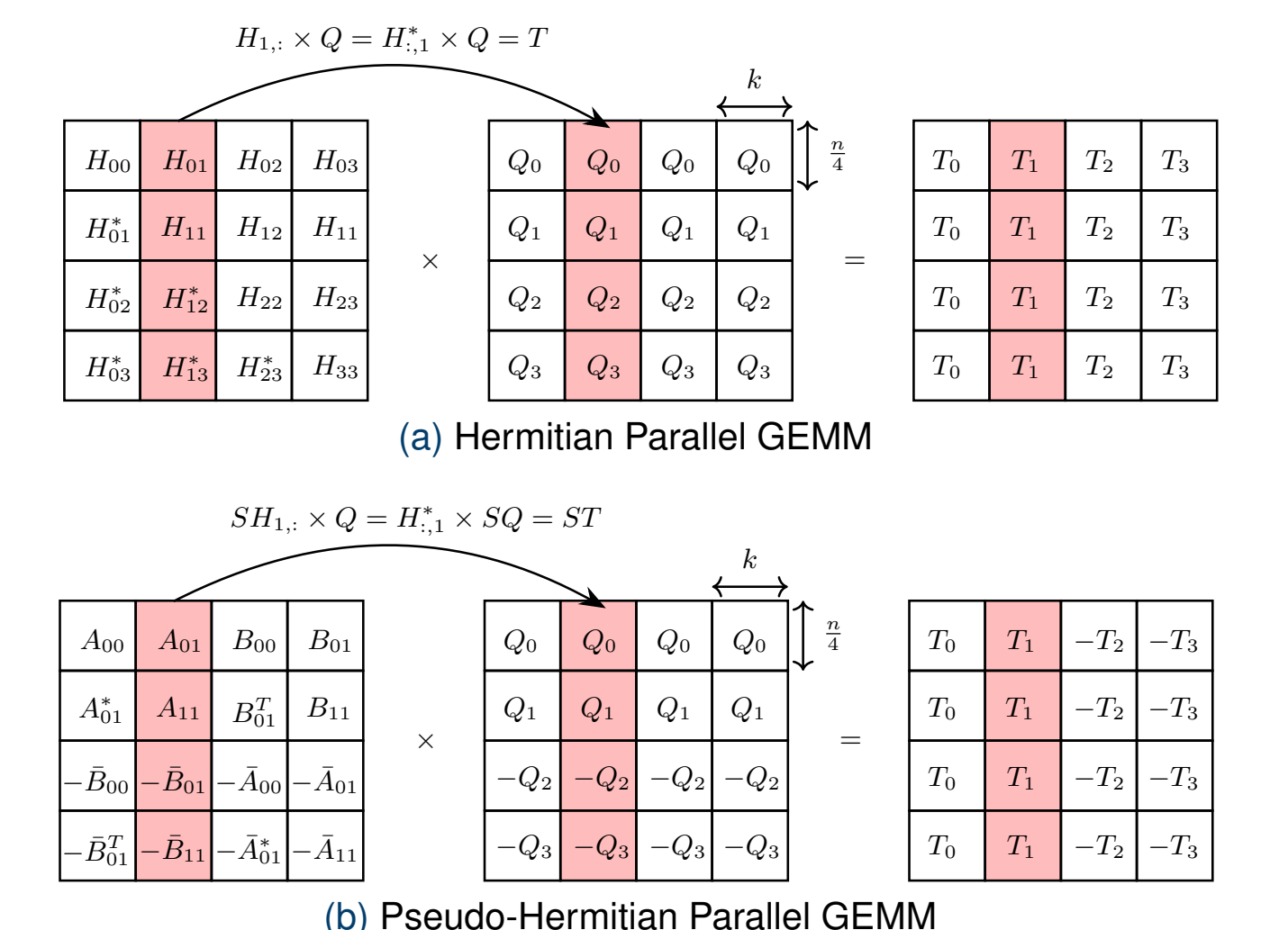


Figure: Flip sign function depending on the data distribution - (gray) sign flipped area, (red) MPI world rank, (blue) MPI column rank

Communication avoiding parallel GEMMs

Chebyshev three-term recurrence relation
 $V_{i+1} = \alpha_i A V_i + \beta_i V_{i-1}, \quad V_1 = A V_0, \quad V_0 = \text{RAND}$



Solving a spectrally equivalent hermitian eigenvalue problem

Setting $M := Q^*SQ$ leads to
 $Q_L := SQ(Q^*SQ)^{-1} \Rightarrow G = (Q^*SQ)^{-1}Q^*SHQ$
which is equivalent to solving the hermitian eigenvalue problem
 $L^*(Q^*SQ)^{-1}Ly = \tilde{\lambda}y, \quad \text{with} \quad y = L^*w \quad \text{and} \quad z = Ly.$
with L the Cholesky factor of the decomposition
 $Q^*SHQ = LL^*$ since SH is HPD.

Construction of the hermitian Rayleigh Quotient with implicit dual space

- 1: Compute $T \leftarrow HQ$ $\triangleright$ BLAS GEMM n^2k
- 2: Flip the sign of the lower half $T \leftarrow ST$ $\triangleright$ INTERNAL FLIP $\frac{n}{2}k$
- 3: Compute $W \leftarrow Q^*T$ $\triangleright$ BLAS GEMM nk^2
- 4: Factorize $L \leftarrow \text{CHOLSKY}(W)$ $\triangleright$ LLAPACK PORTF k^3
- 5: Compute $M \leftarrow I - 2Q_2Q_2^* (= Q^*SQ)$ $\triangleright$ BLAS GEMM $\frac{n}{2}k^2$
- 6: Copy $G \leftarrow M$ $\triangleright$ BLAS COPY k^2
- 7: Solve $G \leftarrow L^{-1}(GL^{-*})$ $\triangleright$ 2 LLAPACK TRSM $2k^2$
- 8: Solve $GY = \Lambda Y$ $\triangleright$ LLAPACK HEEVD k^2
- 9: Back-transform & Return $\tilde{\lambda}, Y \leftarrow \tilde{\lambda}^{-1}, L^{-1}Y$ $\triangleright$ LLAPACK TRSM k^2

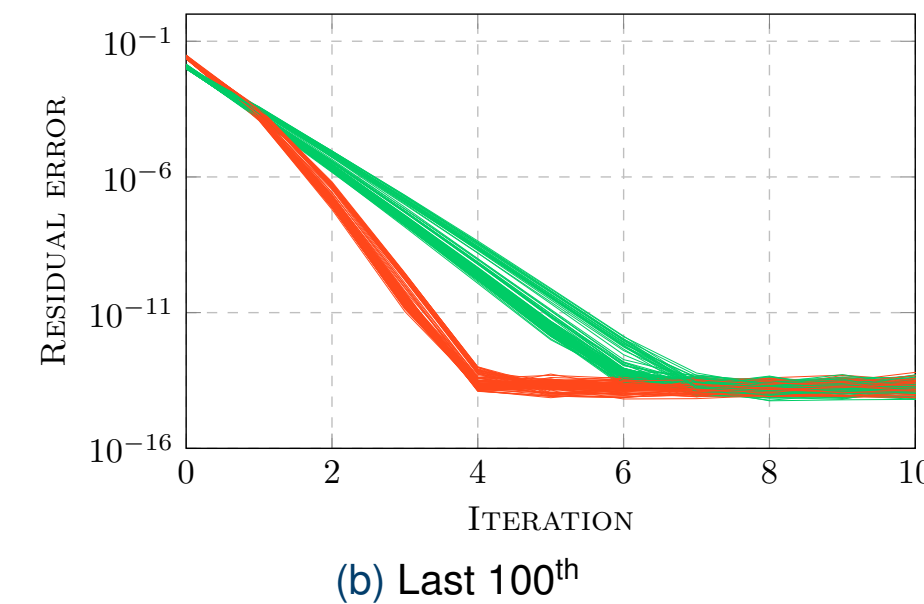
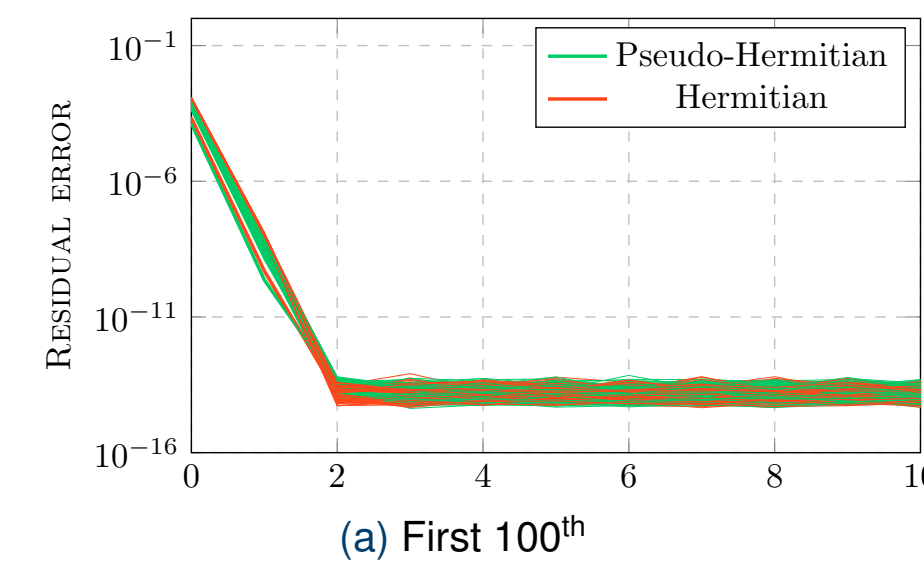


Figure: Evolution of the residuals over $nev = 1000$ for a Silicon pseudo-hermitian H of size $n = 23552$ generated with Yambo

Key benefits of the Hermitian Rayleigh-Quotient

- Equivalent quality for left and right eigenspaces $\|(I - \Pi(Q_L))u\|_2 = \|(I - \Pi(Q))v\|_2 \Rightarrow \mathcal{O}(\sigma_v) = \mathcal{O}(\sigma_u)$
 - Bounded $|\tilde{\delta}|^{-1}$, such that $|\tilde{\delta}|^{-1} \leq \frac{\sqrt{\text{cond}(H)}}{|\lambda_{\min}(Q^*SQ)|}$
- These two features lead to **quadratic convergence** :
 $|\lambda - \tilde{\lambda}| \leq \kappa \cdot \mathcal{O}(\sigma_v^2) \quad \text{with} \quad \kappa := \frac{\sqrt{\text{cond}(H)} \cdot \|H - \lambda I\|_2}{|\lambda_{\min}(Q^*SQ)|}.$

Alternative Rayleigh-Ritz

Setting $M := \text{diag}(Q^*SQ)$ leads to
 $Q_L := [SQ - Q(Q^*SQ - \text{diag}(Q^*SQ))] \text{diag}^{-1}(Q^*SQ),$
and the Rayleigh-Quotient G is given by
 $G = \text{diag}^{-1}(Q^*SQ) [Q^*SHQ - (Q^*SQ - \text{diag}(Q^*SQ))Q^*HQ].$
This alternative Rayleigh-Ritz requires a **non-hermitian eigensolver** (GEEV).

Numerical and Parallel Performance Evaluation

Testing Platforms

Component	JUWELS Booster @JSC	Fugaku @Riken-RCCS
Nodes	936 nodes	158,976 nodes
CPU	2x AMD EPYC 7402, 96 threads	FUJITSU A64FX, 48 cores @ 2.2 GHz
CPU Memory	512 GB DDR4 RAM	32 GB HBM2 @ 1024 GB/s
GPU	4x A100 (40 GB) / NVLink 3	
MPI grid	1 node / 4 GPUs / 4 ranks	1 node / 4 ranks / 48 threads

Convergence

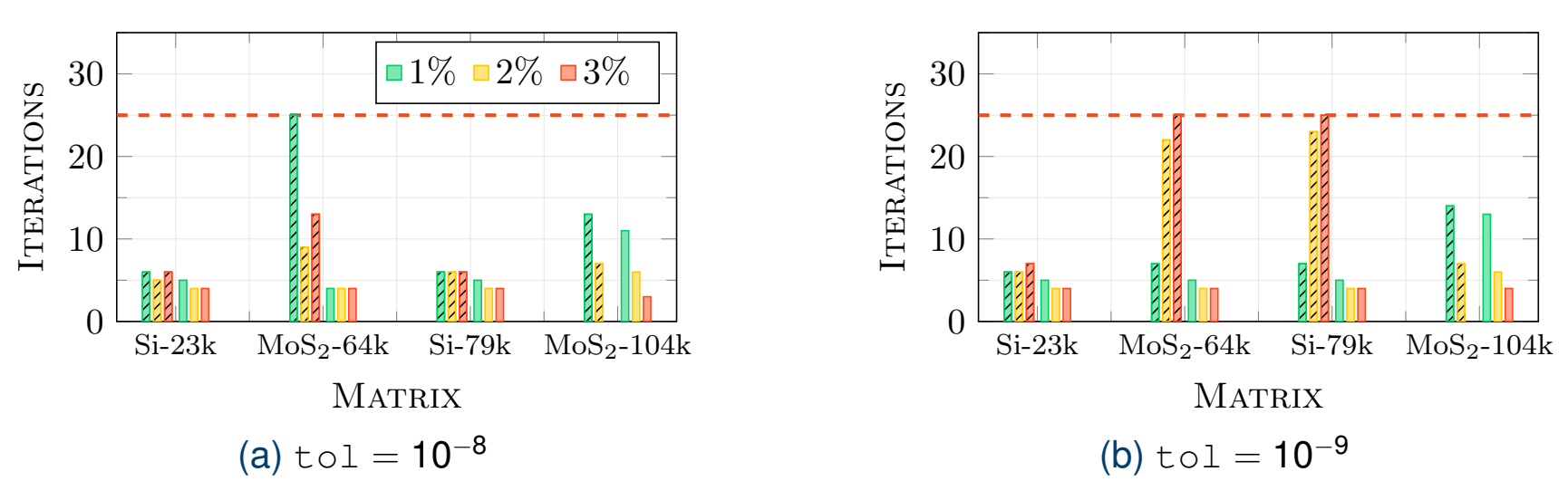


Figure: Number of iterations with respect to the matrix size - GEEV (dashed bars) vs. HEEVD (plain bars)

Percentage of TPP

n	JUWELS Booster - 16 GPUs			Fugaku - 64 ranks small & med vs. 256 ranks large		
	nev = n x 1%	nev = n x 2%	nev = n x 3%	nev = n x 1%	nev = n x 2%	nev = n x 3%
23552	36.2	43.8	48.4	29.872	31.642	32.874
64512	78.2	78.1	78.2	42.9	41.9	41.3
104832	86.8	84.4	81.3	42.2	38.7	33.7

Table: Percentage of TPP occupied by ChASE for each matrix size with respect to nev

Benchmark on JUWELS Booster

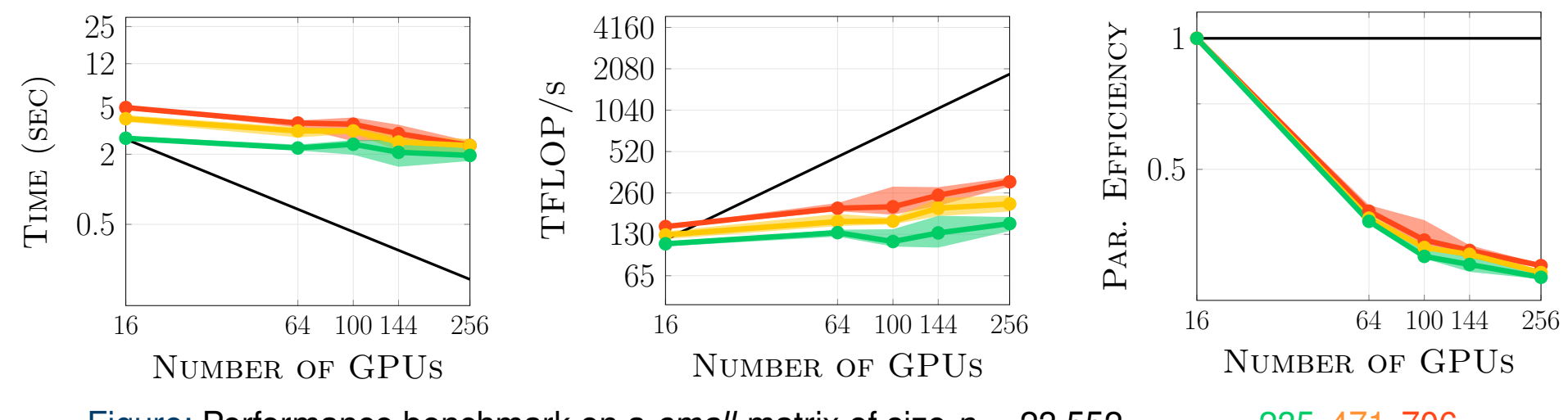


Figure: Performance benchmark on a small matrix of size $n = 23,552$ - $nev = 235,471,706$

Benchmark on Fugaku

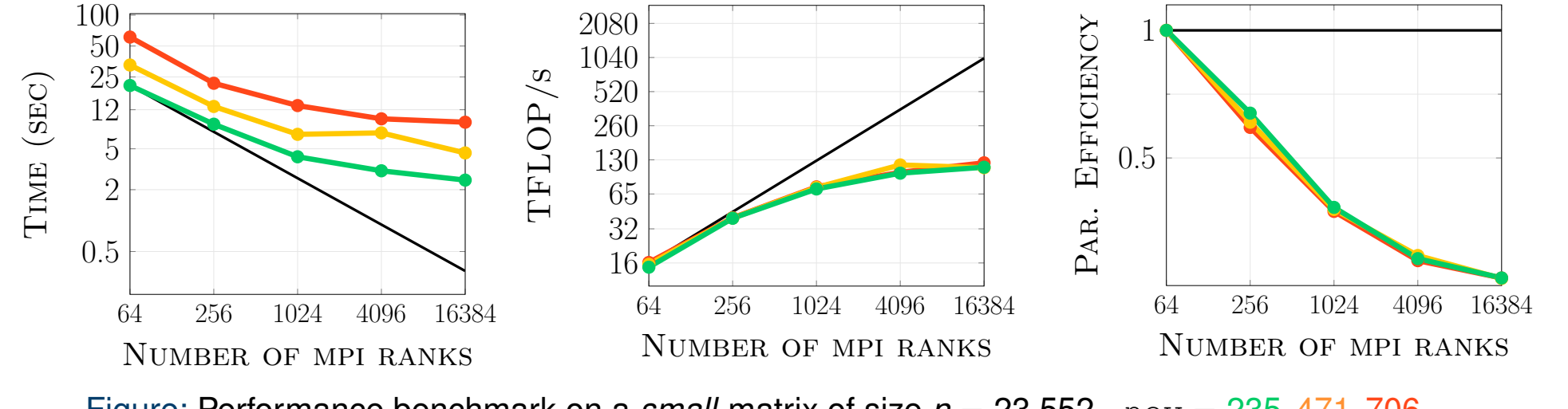


Figure: Performance benchmark on a small matrix of size $n = 23,552$ - $nev = 235,471,706$

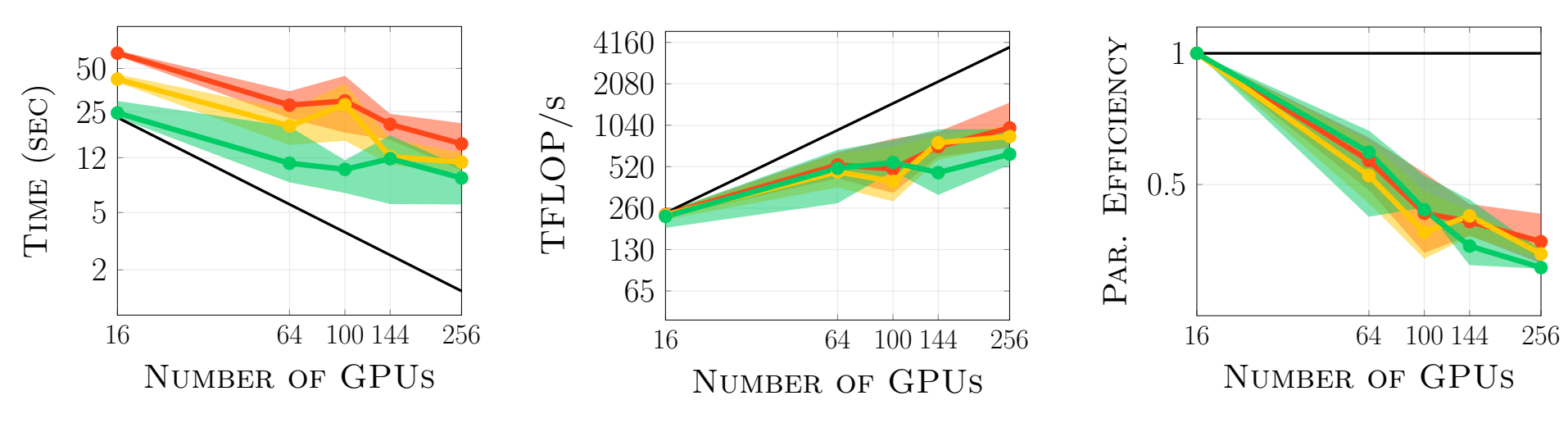


Figure: Performance benchmark on a medium matrix of size $n = 64,512$ - $nev = 645,129,1935$

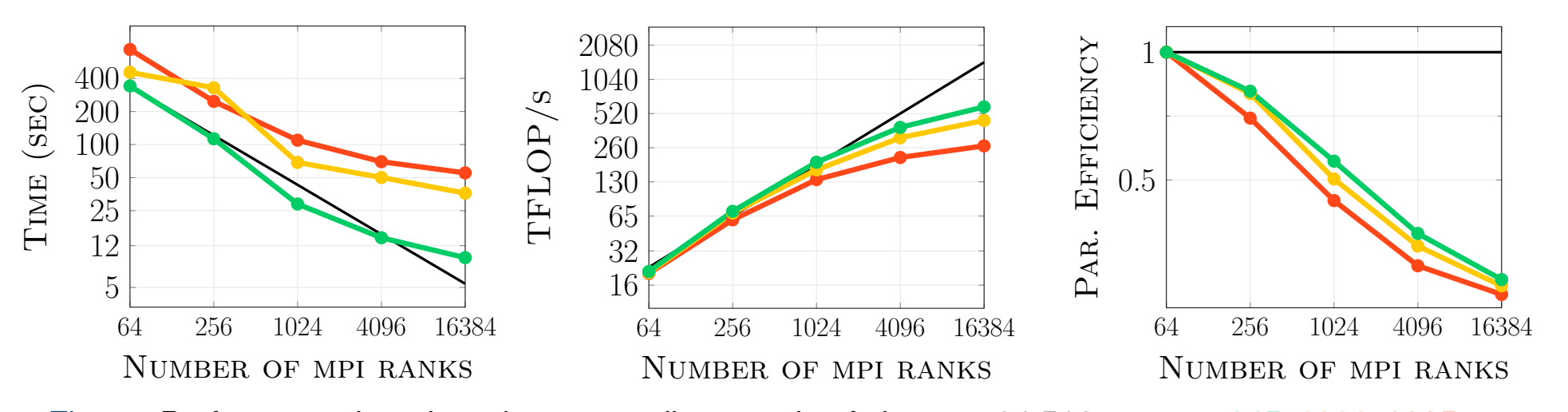


Figure: Performance benchmark on a medium matrix of size $n = 64,512$ - $nev = 645,129,1935$

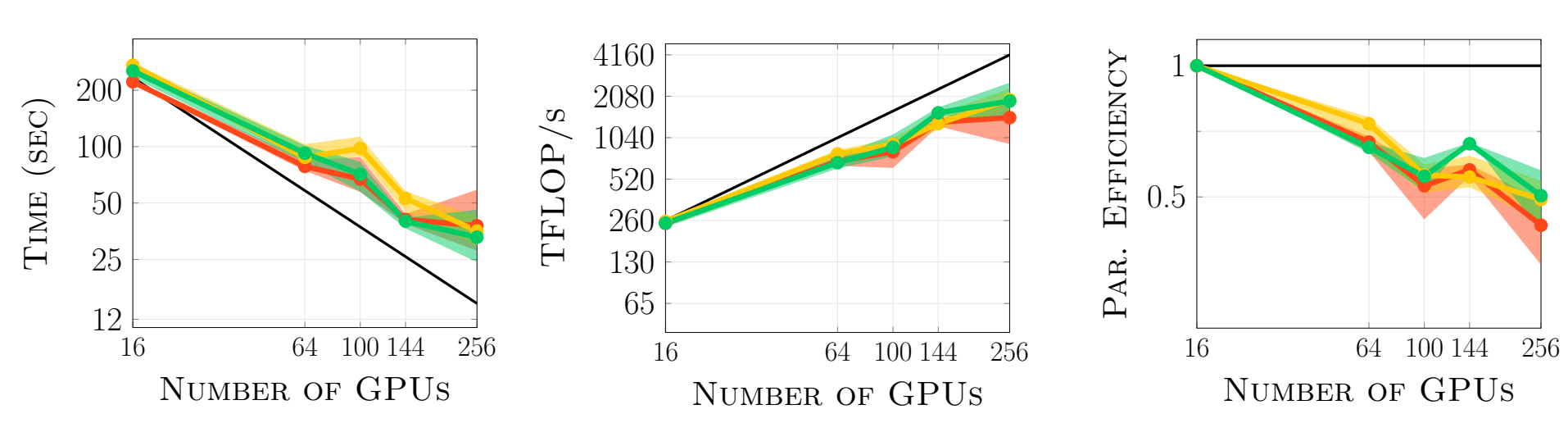


Figure: Performance benchmark on a large matrix of size $n = 104,832$ - $nev = 1048,2096,3144$

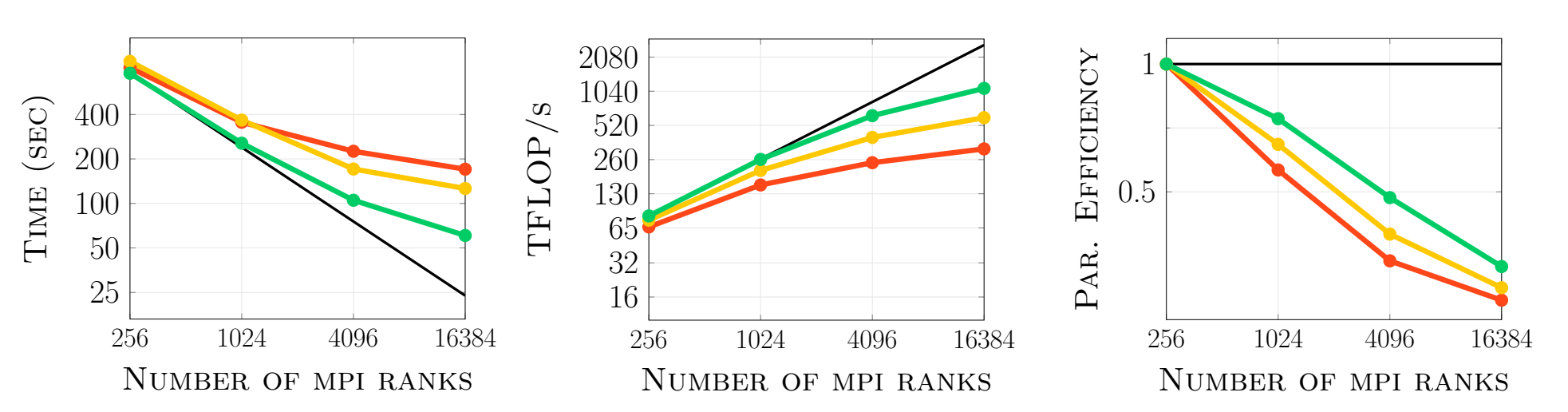


Figure: Performance benchmark on a large matrix of size $n = 104,832$ - $nev = 1048,2096,3144$

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